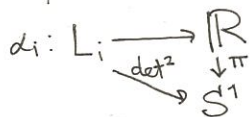


Lecture

Dimension of $\mathcal{M} = \{ (u, S) \}$

Last time Given two graded Lagrangians

$L_i = (L_i, \alpha_i) \quad i=0,1$



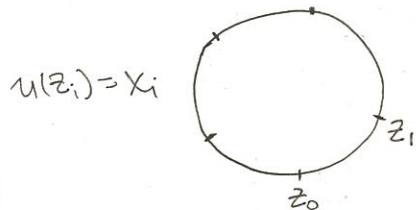
Goal of today: sketch an argument for the following

claim: Fix $(L_i, \alpha_i) \in \mathcal{M} \quad i=0,1,\dots,k$. And $x_i \in L_{i-1} \cap L_i$

Let $(u, S) \in \mathcal{M} \leftarrow$ maps: $u: S \rightarrow M$ w/ $\partial u \subset \bigcup_{i=0}^k L_i$,
marked pts $\rightarrow x_i, (u - X_Y)^{\circ,1} = 0$

Then $\dim \mathcal{M} = i(x_0) - \sum i(x_i) + (k-2)$

↑ incoming ↑ outgoing



$u(z_i) = x_i$

Label z_0 as "incoming" ∂ pt.
 z_1, \dots, z_k as "outgoing" ∂ pt.
Let $i(x_i) :=$ [the integer from the last class]

↑ signed intersection of \tilde{X} with $\alpha_i(x) + \mathbb{Z}$. or "almost winding number"

$x_0 \in CF(L_0, L_k)$

$x_1 \in CF(L_0, L_1)$

$\mu^k: x_k \otimes \dots \otimes x_1 \rightarrow \sum \# \cdot x_0$

i.e. $\mathcal{M}_u \subset \mathcal{M}$ - component containing u

$\dim \mathcal{M}_u = \bigcirc \iff i(x_0) = \sum_{\text{outgoing}} i(x_i) + 2 - k$

this is why μ^k is a map.

$CF(L_{k-1}, L_k) \otimes \dots \otimes CF(L_0, L_1) \rightarrow CF(L_0, L_k) [2-k]$

Def: As a graded ab. gp: $CF(L_0, L_1) = \bigoplus_{x_i} \mathbb{Z} [-i(x_i)]$.

Let's first examine $S = \mathbb{R} \times [0,1]$

Fix a map $u: S \rightarrow M$ satisfying 2 conditions,
pseudoholomorphic $((du - X_Y)^{\circ,1} = 0)$

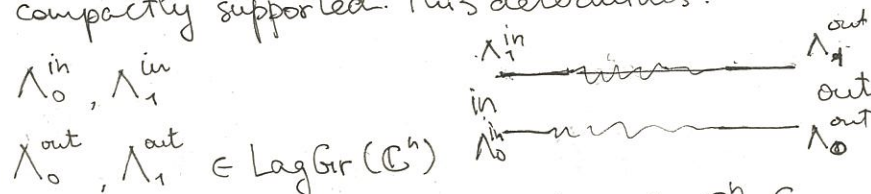
Last time: I claimed you could compare $\dim \mathcal{M}_u$
by examining some winding number $\mu_{\partial S}: \partial S \rightarrow \text{LagGr}$

$\mathbb{R} \times \{1\} \rightarrow \text{LagGr}(u^*TM) \cong S \times \text{LagGr}(\mathbb{C}^n) \rightarrow \text{LagGr}(\mathbb{C}^n)$

 $\mathbb{R} \times \{0\} \rightarrow \text{LagGr}(u^*TM) \cong S \times \text{LagGr}(\mathbb{C}^n) \rightarrow \text{LagGr}(\mathbb{C}^n)$

How to make a loop in LagGr?

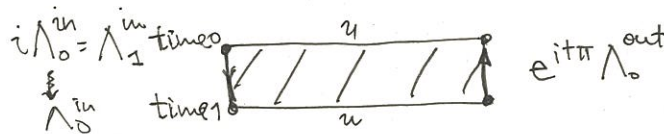
• Assemble the maps $\mathbb{R} \times \{i\} \rightarrow \text{LagGr}(\mathbb{C}^n)$
compactly supported. This determines:



Can find a trivialization of the bundle $\mathbb{C}^n \times S$

$\Lambda_1^{\text{in}} = i \Lambda_0^{\text{in}}$ and $\Lambda_1^{\text{out}} = i \Lambda_0^{\text{out}}$

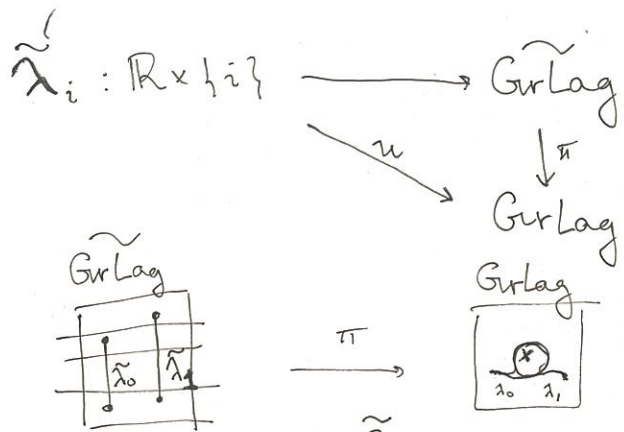
This defines a loop in $\text{LagGr}(\mathbb{C}^n)$:



Thm (Fiber) $\dim \mathcal{M}_u =$ winding # of this path

How to relate to $i(x_{\text{in}}), i(x_{\text{out}})$?

Assume we choose lift:



We can choose path \tilde{c} in $\widetilde{\text{GrLag}}$ connecting the endpoints of $\tilde{\lambda}$.



So computing the bdy #'s given by λ 's is the same as computing winding #'s given by the \tilde{c} .

But the indices $i(x_i)$ were defined to compute the indices of \mathbb{C} .

(The area of the strip //)

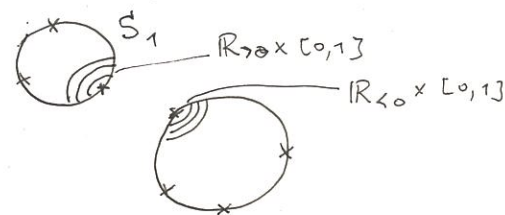
For other cases: "think about" exact same.

"Actual Proof" gluing $\bar{\partial}$ operators over disks

If $S = S_1 \# S_2$

\leftarrow holomorphic connect sum

Then $\text{index}(\bar{\partial}_S) = \text{index}(\bar{\partial}_{S_1}) + \text{index}(\bar{\partial}_{S_2})$



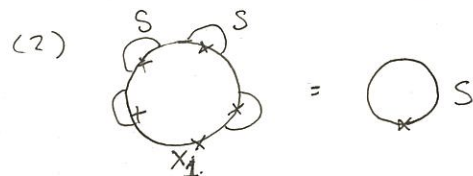
Briefly: When you take the derivative D of $(- - x_i)^{0,1}$ get a Fredholm operator (i.e. $\dim \ker D, \dim \text{coker} D < \infty$)

Turns out ~~deform~~ we can deform D to standard $\bar{\partial}$ through Fredholm operators.

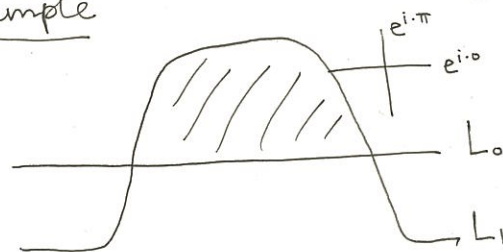
Fact: index doesn't change under Fredholm ~~operator~~ path. When D is surjective (condition) $\text{index} = \dim \ker - \dim \text{coker} = \dim \ker = \dim M_u$ it follows $\text{coker} = 0 \Rightarrow \dim \text{coker} = 0$

Gives the gluing formula:

(1) Can compute index for $S = \mathbb{D}^2 \setminus \text{single pt.}$



Example



WILL COME BACK TO IT NEXT TIME

$M = \mathbb{C}, dx \wedge dy, J = i$

Choose trivialization of $\det^2(TM)$ so that $\text{---} |_{i \in S^1} |_{i \in S^1}$