

19/10 MATH - 297 (TFT and Hochschild)

exercise 1st

$$z: \text{Cob}^{\perp} \rightarrow \text{Vect}_K^{\otimes}$$

be a TFT. Show that $z(\emptyset)$

must be a finite dim vector space.

Today: intro to TFTs.

Defn. An n -dim topo. field theory

is a symmetric monoidal

functor $Z: (\text{Cob}^{\text{fr}})^{\perp} \rightarrow \mathcal{C}^{\otimes}$ (^{partition function})

where \mathcal{C}^{\otimes} is some (∞, n)

- category w/ \otimes as symmetric monoidal
- lecture.

Rmk. There's a theorem proven (Ayala - Francis) that gives an algebraic characterization of all possible \mathcal{Z} .

Thm ("cobordism hypothesis") -
— à la Baez-Dolan,
Lurie ...)

$$\{ \mathcal{Z} \} \simeq \{ \text{redundant fully dualizable objects of } \mathcal{C} \}$$

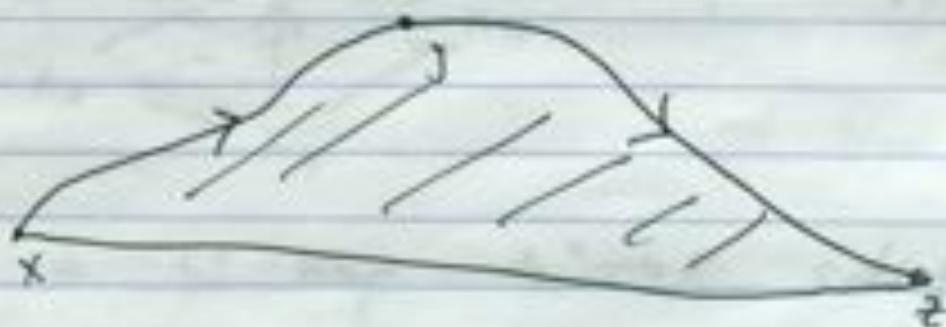
The proof is completely different from Lurie strategy. Chris is trying to follow the original program.

(∞, n) -cat. For your sanity, all definitions are heuristic.

$n=0$. $(\infty, 0)$ -cat = space.

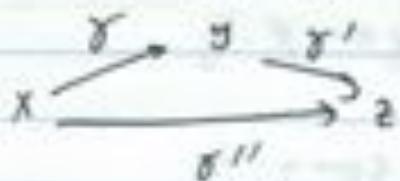
Given a top space X , define a "category" X :

- $ob(X) = \{$ points of top. $\}$
space X
- $hom(x, y) = \{$ cont. $\gamma : [0, 1] \rightarrow X$
 $\gamma(0) = x, \gamma(1) = y$
- Instead of defining composition, we'll



say what a commutative
diagram should be!

A commutative triangle

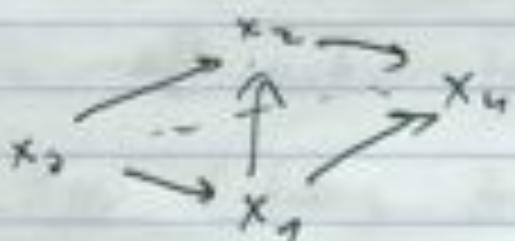


is a co-map $\Delta^n \rightarrow X$

$$\text{s.t. } \partial_0 \Delta^n = \gamma^1, \quad \partial_1 \Delta^n = \gamma^2, \quad \partial_2 \Delta^n = \gamma.$$

A commutative diagram is X is

the form of an n -simplex $\Delta^n \subset R^{n+1}$



$$\left\{ \begin{array}{l} (\alpha_{v_i}, x_{v_i}) \mid \sum x_i = 1 \\ v_i \in \text{vertices} \end{array} \right\}$$

is a co-map $\Delta^n \rightarrow X$ satisfying
the obvious boundary conditions.

(The obvious inclusion $R^n \rightarrow R^{n+1}$, induce

$\Delta^{n-1} \rightarrow \Delta^n$ and $\partial^j \Delta^n \subset \Delta^n$ i.e. delete j^{th} vertex
 $\{ (x_0, \dots, \hat{x_j}, \dots, x_n) \in \Delta^n \mid x_j = 0 \}$ + interior.

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④ Δ Fixing two morphisms

$$\delta: X \rightarrow Y$$

$$\delta': Y \rightarrow Z$$

there exist ∞ -many commutative triangles

$$\begin{array}{ccc} & y & \\ x & \xrightarrow{\quad r \quad} & y' \\ & \xrightarrow{\quad \delta' \quad} & z \end{array}$$

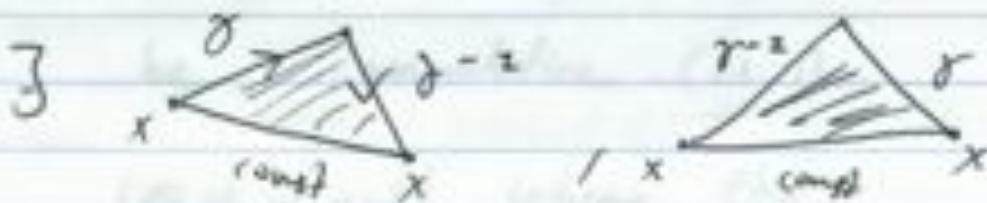
However, we can prove that the space of such triangles is contractible. (follows from weak Kan condition)

Rank every morphism in X is invertible up to homotopy.

In classical category theory (e.g. MacLane)

$$f: X \rightarrow Y \text{ is invertible} \Leftrightarrow \begin{array}{c} f \xrightarrow{\quad \sim \quad} f^{-1} \\ \downarrow id_X \qquad \downarrow id_Y \end{array}, \quad \begin{array}{c} f^{-1} \xrightarrow{\quad \sim \quad} f \\ \downarrow id_X \qquad \downarrow id_Y \end{array}$$

exist ($\exists f^{-1}$ and these are commutative triangles). This in our case is obvious:



For that reason, an $(\infty, 0)$ -cat is an ∞ -groupid. The " \circ " in $(\infty, 0)$ -category is the fact that every k -morphism for $k > 0$ is invertible.

n=2: An $(\infty, 1)$ -category is

- A category where

- $\text{Hom}(x, y)$

- is a topo. space (i.e.

- a category enriched in topo. spaces)

the idea of what it

means for a k -simplex to

be commutative (as in

$(\infty, 0)$ - cat) where the

edges NEED NOT be

invertible.

connected space of possible compositions

is still contractible.

Example. Let X be a space of

stratification: (Jacob's def). a space

X with a continuous map $X \rightarrow P$

where P is a poset. P is

topo \Rightarrow that $u \in P$ open $\Leftrightarrow u$ is open
closed

i.e. if $x \in U$, then any $y \geq x$

is in U . $\Rightarrow U$ is closed.

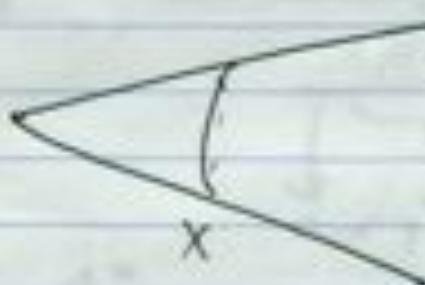
Sub Example. $X = \emptyset, P = \emptyset$

Given $X = \emptyset$ stratified, consider

$$C(X) := \begin{matrix} X \hookrightarrow X \times [0, \infty) \\ \downarrow \quad \downarrow \\ * \longrightarrow C(X) \end{matrix}$$

the open cone. When $X \neq \emptyset$ this is

$$X \times [0, \infty) / X \times \{0\}$$



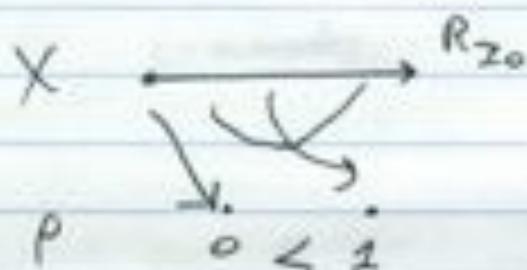
and $C(\emptyset) = \emptyset$ w/ new minimal

element adjoint.

sub-example. ① If $x = p = \emptyset$, then $C(X) = ct$

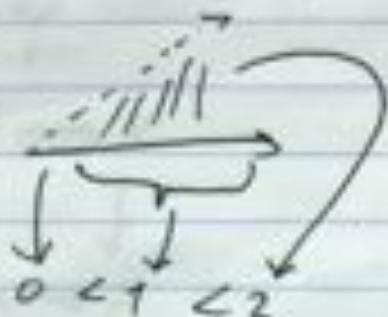
and $C(P) = ct$, so $C(X) = ct$
 \downarrow \downarrow
 $C(ct) = ct$

② $X = P = ct$



Space X is now two pieces.

③ Repeat ...



③ Δ^n
0 < i < n

The example (finally!) is the ex.7
path category of a stratified space X .

objects = points of X

$\text{mor}(x,y) =$ space of continuous
functions $C(X) \cdot R_{\geq 0} \xrightarrow{r} X \xrightarrow{\gamma} Y$

s.t. diagram $\downarrow \quad \downarrow \quad \gamma(x) = y$

$\text{O} \in + \xrightarrow{\rho} P$
 map of
 point
 $\mathcal{C} = \text{continuous} \Leftrightarrow$
 map of points

A comm. diagram in shape of a
is a co-map

$$\begin{array}{ccc} \Delta^k & \longrightarrow & X \\ \downarrow & & \downarrow \\ \text{O} \in \text{co} \dashv & \longrightarrow & P \end{array}$$

Example. $X = \frac{x_0 - z}{z - y_0}$



perfect

Claim. $\ker(X_0, y_0) = \emptyset$. (consider

$$\begin{array}{ccc} R_{z_0} & \xrightarrow{\delta} & X \\ \downarrow & & \downarrow \\ \text{obj} & \longrightarrow & \text{obj} \end{array}$$

②

Also, a path from x_0 to z

is not invertible.

Example. Cob₂ fr. Objects are disjoint union of nifls of the form $gt \times R +$ orient.

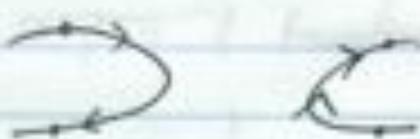
You should ignore the R -component. Think its + infinitesimal ~~assumptions~~ without orientations.

A morphism is an oriented cobordism.

Example. $X = \begin{array}{c} \rightarrow + \\ \leftarrow - \end{array}$ $Y = \begin{array}{c} \rightarrow + \\ \leftarrow - \end{array}$

Some morphism $X \rightarrow Y$:

$$\text{id}: \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} X \times \mathbb{R}$$



Composition is as usual.

Note $\text{Cat}_{\geq 2}$ has a symmetric monoidal

structure given by disjoint union.

n ≥ 2. An (∞, n) -category is a

gadget with

- * objects x, y, \dots

- (1-) * morphism $\text{hom}_{\psi_{f,g}}(x,y), \dots$

- (2-) * morphism between morphism $\text{hom}(f,g)$

- till we have a space of
isomorphisms.

Example: \mathcal{C} = cat of categories

ob: X is a category

$$\text{hom}(X,Y) = \{ \text{functors } f: X \rightarrow Y \}$$

$$\text{hom}(F,G) = \left\{ \begin{array}{c} f \\ \circ \\ g \end{array} \right. \quad \begin{array}{l} \text{the set} \\ \text{of nat.} \\ \text{trans.} \\ y : F \rightarrow G \end{array}$$

Example: \mathcal{C} = Morita category

ob: X is an assoc algebra

$$\text{hom}(A,B) = \left\{ \begin{array}{c} \text{obj: modules} \\ M \end{array} \right\}$$

$$\text{hom}(M,N) = \left\{ \begin{array}{c} \text{bimodules} \\ M \end{array} \right\}$$

M

$X \quad Y$

There is a functor from one to the other called
"forget module".