

last time. Meth 297 (4/11)  
We saw example

of CY-aly where  $\dim \neq 0!$

( $A = \mathbb{R}^i_{\text{DR}}(X)$ ,  $X$  compact or

$\text{tr}: A \rightarrow k[t, d]$ ,  $\alpha \mapsto \int \langle \alpha \rangle$   
[X]

Q. Can we define a version  
of  $D$  that in operator degree  
shift? If so, call it  $D^{\text{op}}$ .

Def. A CY-aly of  $\dim = d$  is  
symm @ functor  $D^d \rightarrow \text{Chains}$ .

A: give well chosen local systems  
to each man-space.

Def. Fix  $X$  a top-space. A local  
system (of cochain cplx) is a  
functor  $X \rightarrow \text{Chains}$ .

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Digression. To make you happy,  
short intro to  $os$ -cat. (A particular  
model of  $(\omega, 1)$ -categories).

Def. Let  $\Delta$  be the cat w/  
objects = finite, linearly ordered  
non-empty sets

Ex.

$$\text{ob } \Delta = \{ [0], [1], [2], \dots \}$$

$\uparrow \quad \quad \uparrow \quad \quad \uparrow$

{0}      {0,1}      {0,1,2}

$$\text{hom}(\mathcal{I}, \mathcal{J}) \stackrel{\text{set}}{=} \left\{ \begin{array}{l} \text{poset map} \\ \mathcal{I} \rightarrow \mathcal{J} \end{array} \right\}$$

, i.e., let  $f: \mathcal{I} \rightarrow \mathcal{J}$  s.t.  $i \leq i'$   
iff  $f(i) \leq f(i')$ . That is just  
weakly order preserving maps.



Def. A simplicial set is a functor

$$\Delta^{op} \rightarrow \text{Sets}$$
$$[i] \mapsto X_i$$

~~Let  $X$  be a topological space~~

Ex. Fix  $n \geq 0$ . We consider

$$\Delta^n := \text{hom}_{\Delta}(-, [n]) : \Delta^{op} \rightarrow \text{Sets}$$

that's a simplicial set!

Ex. Fix  $X$  a space. Then

define

$$\Delta^{op} \rightarrow \text{Sets}$$
$$[n] \mapsto \text{Hom}_{\text{spaces}}(\Delta^n, X)$$

↑  
space

i.e. continuous functions

What does it do to morphisms?

Note  $\exists$  functor

$\Delta \rightarrow \text{Spaces}$

$[i] \mapsto \Delta^i =$

$$\left\{ (t_0, \dots, t_i) \mid \sum_{j=0}^i t_j = 1, t_j \geq 0 \right\}$$

$$\uparrow \\ \mathbb{R}^{i+1}$$

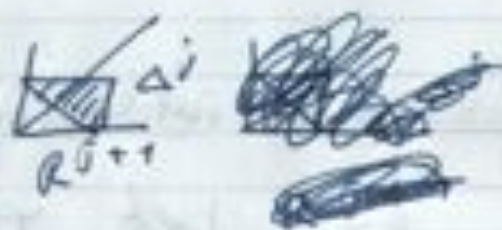
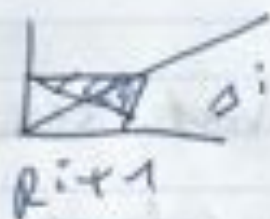
Given a poset map  $\mathcal{P} \rightarrow \mathcal{Q}$ , there is

an induced map  $\mathbb{R}_{\Delta^{\mathcal{P}}} \rightarrow \mathbb{R}_{\Delta^{\mathcal{Q}}}$

Ex. If  $\mathcal{P} = [n]$  and  $\mathcal{Q} = [n+1]$  where  $s_{\mathcal{P}} = \sum_{i \in \mathcal{P} - \{j\}} t_i$

then  $\mathbb{R}_{\Delta^{\mathcal{P}}} \xrightarrow{\theta_j} \mathbb{R}_{\Delta^{\mathcal{Q}}}$  is "skip an element  $j \in \mathcal{P}$ ". What's going on? Pictorially





This induces the inclusion

$$\Delta^i \hookrightarrow \Delta^{i+1}$$

of the \$j^{\text{th}}\$ face.

So compose with the functor

$$\Delta^r \xrightarrow{\Delta} \text{Spaces}^r \xrightarrow{\text{hom}(-, \Delta)} \text{Set}$$

You've seen this before!

Ex.  $C_*^{\text{sing}}(X)$  is obtained

by taking a free abelian group generated by each  $x_i$ ,

then defining the \$i^{\text{th}}\$ differential as the alternating sum of \$d\_j\$'s.

More interesting and relevant example.

Let  $\mathcal{C}$  be a category.

This defines a functor

$$\begin{aligned} N(\mathcal{C}) &\xrightarrow{\quad} \text{Set} \\ [i] &\mapsto \text{fun}([i], \mathcal{C}) \end{aligned}$$

Note that any poset is a category.

Elements of poset =  
obj of category

$$\text{hom}(\mathcal{C}, \mathcal{C}) = \begin{cases} * & , p \leq q \\ \emptyset & , \text{else} \end{cases}$$

Then

$$N(\mathcal{C})_0 = \text{fun}([0], \mathcal{C}) \cong_{\text{Set}} \text{ob } \mathcal{C}$$

cat w/  
one object  $0$ ,  
all unique homs  
 $0 \rightarrow 0$



Idea.

every category gives an example  
of simplicial set ...

$$N(\mathcal{C})_1 = \text{Fun}([1], \mathcal{C})$$

cat w/ obj "0", "1"  
and  $\begin{matrix} \text{id} \\ \downarrow \\ \text{0} \end{matrix} \rightarrow \begin{matrix} \text{id} \\ \downarrow \\ \text{1} \end{matrix}$

$$\cong \coprod_{\substack{\text{set } x_0, x_1 \\ \in \text{obj}(\mathcal{C})}} \text{hom}_{\mathcal{C}}(x_0, x_1)$$

What about morphisms? consider

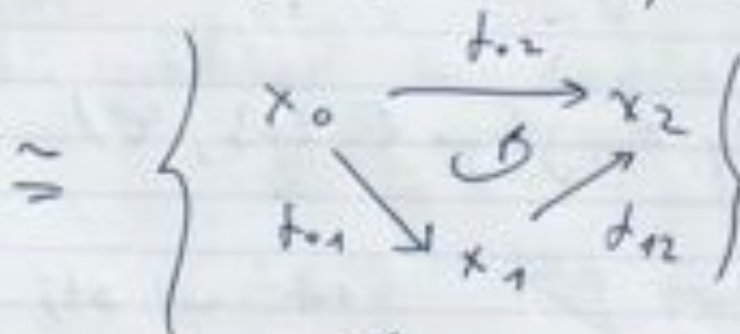
$$[0] \begin{matrix} \xrightarrow{d_0} \\ \xrightarrow{d_1} \end{matrix} [1]$$

$$\begin{cases} d_0(0) = 1 \\ d_1(0) = 0 \end{cases}$$

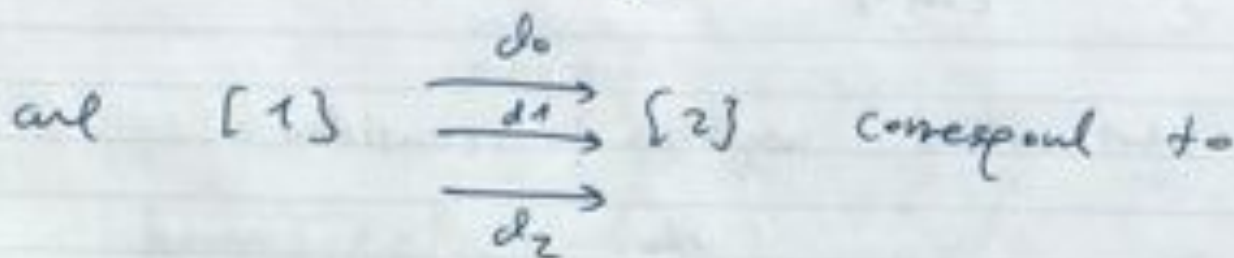
give  $\text{obj}(\mathcal{C}) = N(\mathcal{C})_0$   $\begin{matrix} \xleftarrow{d_0 = \text{target}} \\ \xleftarrow{d_1 = \text{source}} \end{matrix} N(\mathcal{C})_1 = \text{Morphisms}(\mathcal{C})$

Let's figure

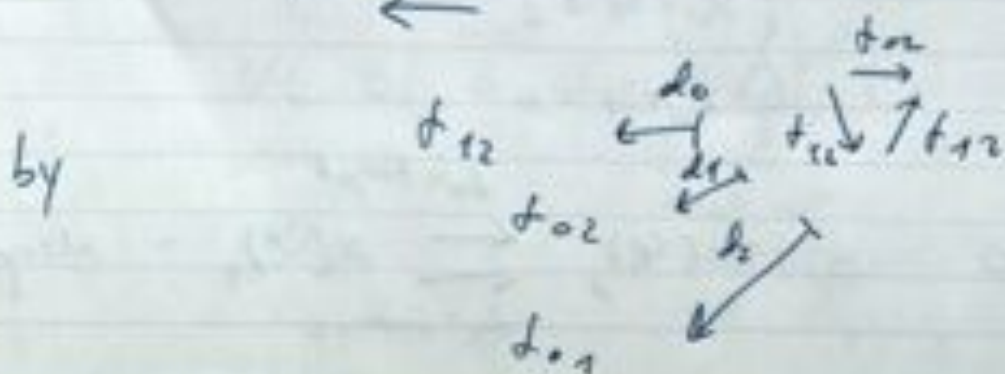
$$N(\mathcal{C})_2 = \text{Fib}([2], \mathcal{C})$$



commutative triangle at  $\mathcal{C}$

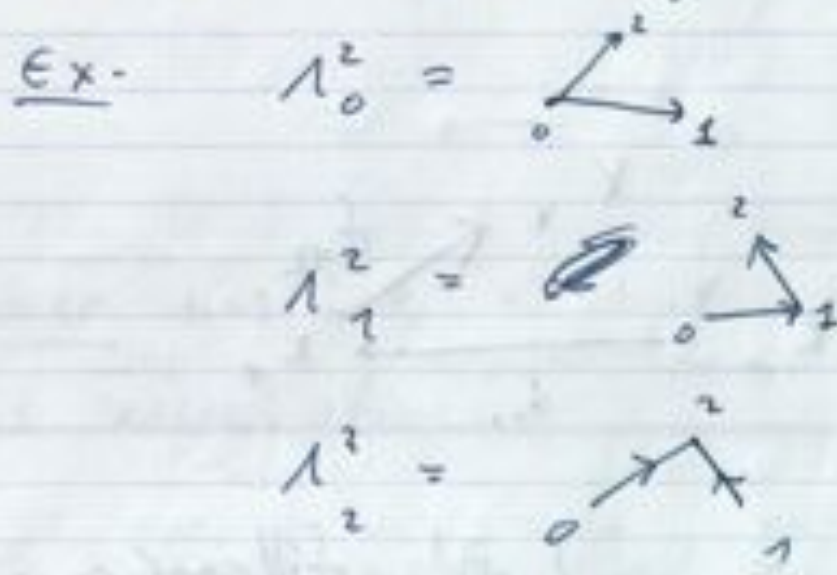


3 maps  $N(\mathcal{C})_1 \leftarrow N(\mathcal{C})_2$  where they are

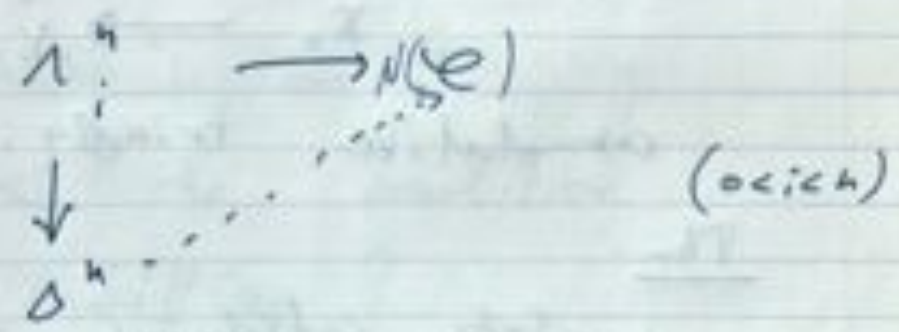




Defn. Let  $\sigma \in \Lambda_K^n \subset \Delta^n$ ,  $0 \leq k \leq n$   
 be the set obtained by removing  
 the  $k^{\text{th}}$  face (opposite vertex  $k$ )



Prop. Fix  $\mathcal{C}$ . Then for every  
 diagram



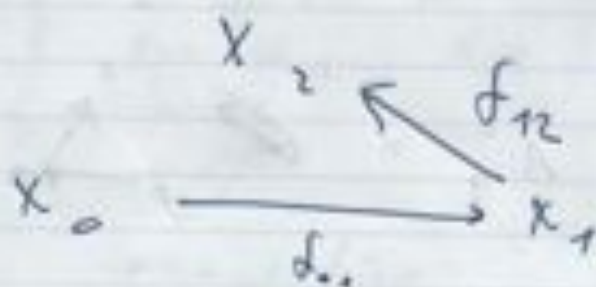
$\exists!$   $f$  making this diagram  
 commute.

Ex.  $n \geq 2, i = 2$ . A functor

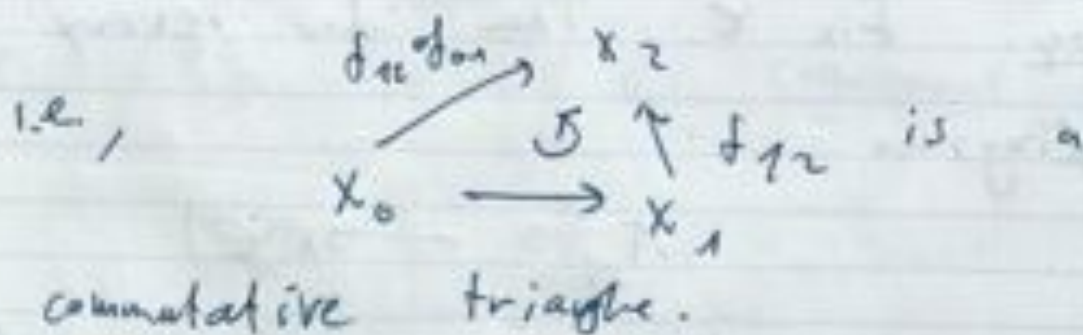
$$\mathcal{A}_1^2 \rightarrow \mathcal{C}$$

is a choice of

#



}  $\exists!$  filler (composition)



Thm.

$\left. \begin{array}{l} s \text{ Sets satisfying} \\ \text{our prop} \end{array} \right\} \begin{array}{l} \cong \\ \text{equiv} \\ \text{of} \\ \text{cat} \end{array} \left\{ \begin{array}{l} \text{small} \\ \text{categories} \end{array} \right.$   
 $\cap$  subcat  
 $s$  Sets



map:

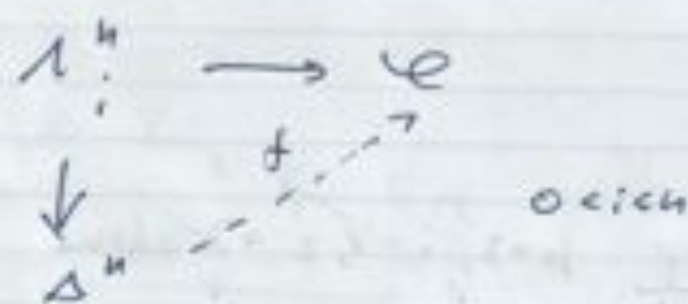
$$N(\mathcal{C}) \longleftrightarrow \mathcal{C}$$

$$x \mapsto \mathcal{C}_x := \left\{ \begin{array}{l} ob = x_0 \\ h = (x, y) \in \mathcal{C}_x \\ \text{whose face} \\ \text{maps to} \\ x, y \end{array} \right.$$

Exer  $k=3, i=1, 2 \Rightarrow$  shows composition  
is associative.

So the condition of filling  
inner horns looks a lot like  
the property of being a  
category. What happens if  
we remove the condition  
of unique filling?

Def. A simplicial set  $\mathcal{C}$  is called an  $\infty$ -cat if  $\forall$  diagram



$\exists \dagger$  making diagram commute.

~~Thm~~ Why nice?

(1) Thm. If  $\mathcal{C}, \mathcal{D}$  are  $\infty$ -cat, then

$$\text{Fun}(\mathcal{C}, \mathcal{D}) := \text{Hom}_{\text{sets}}(\mathcal{C}, \mathcal{D})$$

(natural trans)

this is an  $\infty$ -cat.



The RMS

hom sets  $(\mathcal{C}, \mathcal{D})$

is a set.

$\text{hom}(\mathcal{C}, \mathcal{D})_i = \text{NatHom}(\mathcal{C} \times \underline{\Delta}^i, \mathcal{D})$

This existed like 20 years before

Joyal said "let's use it to

define categories". Very classical.

In almost every other world, this  
~~that~~ <sup>That is</sup> very hard to prove

(2) Great Defn of (homotopy)

(co) limits.

(3) Amazing untheoretical constructions

Probably the best thing in ex-act!

Classical. Given a functor

$\mathcal{C} \xrightarrow{f} \text{sets}$

$\exists (! \text{ up to equiv})$  a category

$\mathcal{E}_f$

$\downarrow$

$\mathcal{C}$

this is kernel of "classifying space construction". Think

$G$ -torsors. Lurie MTT calls it

"(un)straightening".

Other models. • complete Segal spaces  
(just in name) ("simplicial spaces")

• Top-enriched cat

• Cat enriched in simplicial sets

• simplicial model cat



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The ~~best~~ <sup>best</sup> thing ~~about~~ <sup>about</sup> Segal  
 spaces is the generalization  
 to  $(\infty, n)$ , which is immediate  
 and hard/impossible for  $(\infty, n)$ -cat  
 in Lurie sense.

One but thing <sup>about the Lurie model ...</sup> - where id come from?  
 the unit

$$N(\mathcal{C})_1 \leftarrow N(\mathcal{C})_0$$

$$\text{id}_X \leftrightarrow X$$

collapse  
all

$$[1] \longrightarrow [0]$$

$$\mathcal{C}_1 \leftarrow \mathcal{C}_0$$

in Lurie's model we have to specify unit  
 for each  $x \in \text{ob}(\mathcal{C}) = \mathcal{C}_0$

The fact you have to choose is  
the only short coming of this  
approach.